

Programming Languages: Functional Programming

Practicals 1: Functions and Definitions

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You should have installed GHC, with its commandline interface GHCi. Open your favourite text editor, create a new plain text file. The filename extension must end in `.hs`. This will be your working file for this practical. Type `ghci <filename>.hs` in the command line to load the working file into GHCi.

1. Define a function `myeven :: Int → Bool` that determines whether the input is an even number. You may use the following functions:

$$\begin{aligned} \text{mod} &:: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int} , \\ (==) &:: \text{Int} \rightarrow \text{Int} \rightarrow \text{Bool} . \end{aligned}$$

(Types of the functions written above are not in their most general form.)

Solution:

$$\begin{aligned} \text{myeven} &:: \text{Int} \rightarrow \text{Bool} \\ \text{myeven } x &= x \text{ 'mod' } 2 == 0 . \end{aligned}$$

2. Define a function that computes the area of a circle with given radius r (using $22 / 7$ as an approximation to π). The return type of the function might be `Float`.

Solution:

$$\begin{aligned} \text{area} &:: \text{Float} \rightarrow \text{Float} \\ \text{area } r &= (22 / 7) \times r \times r . \end{aligned}$$

3. Part-time students in Institute of Information Science are paid NTD 130 per hour. Define a function `payment :: Int → Int` that, when applied to the numbers of weeks a student work, compute the amount of money the Institute has to pay the student.

- (a) Assume that there are five working days in a week, eight working hours per day. Define *payment*. For clarity, use **let** to define local variables recording number of days worked, number of hours worked, etc.

Solution:

```
payment :: Int → Int
payment weeks = let days = 5 × weeks
                  hours = 8 × days
                  in 130 × hours .
```

- (b) Define *payment* again, but declare the local variables using **where**. Which style do you prefer?

Solution:

```
payment :: Int → Int
payment weeks = 130 × hours
  where hours = 8 × days
        days = 5 × weeks .
```

- (c) The regulation states that students are considered workers, and if a worker works for more than 19 weeks, the Institute has to pay, in addition to the salary, health insurance and pension reserves for the worker. The amount is 6% of the worker's salary.

Update definition of *payment* in the form:

```
payment :: Int → Int
payment weeks | weeks > 19 = ...
              | otherwise = ...
```

You may need a function *fromIntegral* to convert Int to Float, and a function *round* that rounds a floating point number to the nearest integer.

In this case, should you use **let** or **where**?

Solution:

```
payment :: Int → Int
payment weeks | weeks > 19 = round (fromIntegral baseSalary × 1.06)
              | otherwise = baseSalary
  where baseSalary = 130 × hours
        hours = 8 × days
        days = 5 × weeks .
```

For this situation, **where** works better than **let**, since we want the scope of *baseSalary* to extend to both guarded branches.

4. More on **let**.

- (a) Guess what the value of *nested* would be. Type it into your working file and evaluated in in GHCi to see whether you guessed right. Note that indentation matters.

```
nested :: Int
nested = let x = 3
         in (let x = 5
             in x + x) + x .
```

Solution: *nested* evaluates to 13, since the *x* in *x + x* refers to 5 and the *x* in *.. + x* refers to 3.

- (b) Guess what the value of *recursive* would be. Try it in GHCi.

```
recursive :: Int
recursive = let x = 3
           in let x = x + 1
               in x .
```

Solution: The computation does not terminate, since the *x* in *x + 1* refers to itself.

5. Type in the definition of *smaller* into your working file.

```
smaller :: Int → Int → Int
smaller x y = if x ≤ y then x else y .
```

Then try the following:

- (a) In GHCi, type `:t smaller` to see the type of *smaller*.
 - (b) Try applying it to some arguments, e.g. *smaller* 3 4, *smaller* 3 1.
 - (c) Use `:t` to see the type of *smaller* 3 4.
 - (d) Use `:t` to see the type of *smaller* 3.
 - (e) In your working file, define a new function *st3* = *smaller* 3.
 - (f) Find out the type of *st3* in GHCi. Try *st3* 4, *st3* 1. Explain the results you see.
6. More practice on curried functions.

- (a) Define a function *poly* such that $\text{poly } a \ b \ c \ x = a \times x^2 + b \times x + c$. All the inputs and the result are of type *Float*.
- (b) Reuse *poly* to define a function *poly1* such that $\text{poly1 } x = x^2 + 2 \times x + 1$.
- (c) Reuse *poly* to define a function *poly2* such that $\text{poly2 } a \ b \ c = a \times 2^2 + b \times 2 + c$.

Solution:

```
poly :: Float → Float → Float → Float → Float
poly a b c x = a × x × x + b × x + c

poly1 :: Float → Float
poly1 = poly 1 2 1

poly :: Float → Float → Float → Float
poly2 a b c = poly a b c 2
```

7. Type in the definition of *square* in your working file.

- (a) Define a function *quad* :: *Int* → *Int* such that *quad* *x* computes x^4 .

Solution:

```
quad :: Int → Int
quad x = square (square x) .
```

- (b) Type in this definition into your working file. Describe, in words, what this function does.

```
twice :: (a → a) → (a → a)
twice f x = f (f x) .
```

- (c) Define *quad* using *twice*.

Solution:

```
quad :: Int → Int
quad = twice square .
```

8. Replace the previous *twice* with this definition:

```
twice :: (a → a) → (a → a)
twice f = f · f .
```

- (a) Does *quad* still behave the same?
- (b) Explain in words what this operator ($\cdot$) does.

9. Functions as arguments, and a quick practice on sectioning.

- (a) Type in the following definition to your working file, without giving the type.

$$\text{forktimes } f \ g \ x = f \ x \times g \ x .$$

Use t in GHCi to find out the type of *forktimes*. You will end up getting a complex type which, for now, can be seen as equivalent to

$$(t \rightarrow \text{Int}) \rightarrow (t \rightarrow \text{Int}) \rightarrow t \rightarrow \text{Int} .$$

Can you explain this type?

- (b) Define a function that, given input x , use *forktimes* to compute $x^2 + 3 \times x + 2$. **Hint:** $x^2 + 3 \times x + 2 = (x + 1) \times (x + 2)$.

Solution:

```
compute :: Int → Int
compute = forktimes (+1) (+2) .
```

- (c) Type in the following definition into your working file: $\text{lift2 } h \ f \ g \ x = h \ (f \ x) \ (g \ x)$. Find out the type of *lift2*. Can you explain its type?

Solution:

$$\text{lift2} :: (a \rightarrow b \rightarrow c) \rightarrow (d \rightarrow a) \rightarrow (d \rightarrow b) \rightarrow d \rightarrow c .$$

- (d) Use *lift2* to compute $x^2 + 3 \times x + 2$.

Solution:

```
compute :: Int → Int
compute = lift2 (×) (+1) (+2) .
```

10. Let the following identifiers have type:

```
f :: Int → Char
g :: Int → Char → Int
h :: (Char → Int) → Int → Int
x :: Int
y :: Int
c :: Char
```

Which of the following expressions are type correct?

1. $(g \cdot f) x c$
2. $(g x \cdot f) y$
3. $(h \cdot g) x y$
4. $(h \cdot g x) c$
5. $h \cdot g x c$

You may type the expressions into Haskell and see whether they type check. To define f , for example, include the following in your working file:

```
 $f :: Int \rightarrow Char$   
 $f = undefined$ 
```

However, it is better if you can explain why the answers are as they are.

Solution:

1. $(g \cdot f) x c$. We calculate:

$$\begin{aligned} & (g \cdot f) x c \\ = & \{ \text{function application binds to the left} \} \\ & ((g \cdot f) x) c \\ = & \{ \text{definition of } (\cdot) \} \\ & (g (f x)) c . \end{aligned}$$

One can then see that there is a type error: $f x$ is a Char, while g expects Int as its first argument.

2. $(g x \cdot f) y$. We calculate:

$$\begin{aligned} & (g x \cdot f) y \\ = & \{ \text{definition of } (\cdot) \} \\ & g x (f y) , \end{aligned}$$

which type-checks because

- we have $g :: Int \rightarrow Char \rightarrow Int$, and
- x is an Int, thus $g x :: Char \rightarrow Int$.
- Furthermore, $f y$ is a Char. Thus $g x (f y) :: Int$.

The result type is Int.

3. $(h \cdot g) x y$. We calculate:

$$\begin{aligned} & (h \cdot g) x y \\ = & \{ \text{function application binds to the left} \} \\ & ((h \cdot g) x) y \\ = & \{ \text{definition of } (\cdot) \} \\ & (h (g x)) y . \end{aligned}$$

Now we reason:

- Recall $h :: (\text{Char} \rightarrow \text{Int}) \rightarrow \text{Int} \rightarrow \text{Int}$.
- Since $g :: \text{Int} \rightarrow \text{Char} \rightarrow \text{Int}$ and $x :: \text{Int}$, we have $g x :: \text{Char} \rightarrow \text{Int}$
- Thus $h (g x) :: \text{Int} \rightarrow \text{Int}$.
- Since $y :: \text{Int}$, we have $(h (g x)) y :: \text{Int}$.

Thus the expression type-checks, with type Int .

4. $(h \cdot g x) c$. We calculate:

$$\begin{aligned} & (h \cdot g x) c \\ = & \{ \text{definition of } (\cdot) \} \\ & h (g x c) . \end{aligned}$$

We reason:

- The part $g x c$ type-checks, since $g :: \text{Int} \rightarrow \text{Char} \rightarrow \text{Int}$, $x :: \text{Int}$, and $c :: \text{Char}$. We have that $g x c :: \text{Int}$.
 - However, $h :: (\text{Char} \rightarrow \text{Int}) \rightarrow \text{Int} \rightarrow \text{Int}$ expects a function of type $\text{Char} \rightarrow \text{Int}$ as an argument, not Int . Thus the expression fails to type-check.
5. $h \cdot g x c$. Similar to the reasoning above, $g x c :: \text{Int}$. However, function composition $(\cdot)$ expect to compose functions together, and $g x c$ is not a function. Thus the expression fails to type-check.