

PROGRAMMING LANGUAGES: FUNCTIONAL PROGRAMMING

1. INTRODUCTION TO HASKELL: VALUE, FUNCTIONS, AND TYPES

Shin-Cheng Mu

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National Taiwan University and Academia Sinica

A QUICK INTRODUCTION TO HASKELL

- We will mostly learn some syntactical issues, but there are some important messages too.
- Most of the materials today are adapted from the book *Introduction to Functional Programming using Haskell* by Richard Bird. Prentice Hall 1998.
- References to more Haskell materials are on the course homepage.

COURSE MATERIALS AND TOOLS

- Course homepage: <https://scmu.github.io/plfp/>
 - Announcements, slides, assignments, additional materials, etc.
- We will be using the Glasgow Haskell Compiler (GHC).
 - A Haskell compiler written in Haskell, with an interpreter that both interprets and runs compiled code.
 - See the course homepage for instructions for installation and other info.

FUNCTION DEFINITION

- A function definition consists of a type declaration, and the definition of its body:

square :: *Int* → *Int*

square x = $x \times x$

smaller :: *Int* → *Int* → *Int*

smaller x y = **if** $x \leq y$ **then** *x* **else** *y*

- The GHCi interpreter evaluates expressions in the loaded context:

? *square* 3768

14197824

? *square* (*smaller* 5 (3 + 4))

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VALUES AND EVALUATION

EVALUATION

One possible sequence of evaluating (simplifying, or reducing)
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EVALUATION

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EVALUATION

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ANOTHER EVALUATION SEQUENCE

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$$\begin{aligned} & \text{square } (3 + 4) \\ = & \quad \{ \text{definition of } \textit{square} \} \\ & (3 + 4) \times (3 + 4) \end{aligned}$$

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ANOTHER EVALUATION SEQUENCE

- Another possible reduction sequence:

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- In this sequence the rule for *square* is applied first. The final result stays the same.
- Do different evaluations orders always yield the same thing?

A NON-TERMINATING REDUCTION

- Consider the following program:

three $:: \text{Int} \rightarrow \text{Int}$

three $x = 3$

infinity $:: \text{Int}$

infinity $= \text{infinity} + 1$

- Try evaluating *three infinity*. If we simplify *infinity* first:

three infinity

$= \{ \text{definition of } \textit{infinity} \}$

three (*infinity* + 1)

$= \textit{three} ((\textit{infinity} + 1) + 1) \dots$

- If we start with simplifying *three*:

three infinity

$= \{ \text{definition of } \textit{three} \}$

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EVALUATION ORDER

- There can be many other evaluation orders. As we have seen, some terminates while some do not.
- *normal form*: an expression that cannot be reduced anymore.
 - 49 is in normal form, while 7×7 is not.
 - Some expressions do not have a normal form. E.g. *infinity*.
- A corollary of the Church–Rosser theorem: an expression has at most one normal form.
 - If two evaluation sequences both terminate, they reach the same normal form.

EVALUATION ORDER

- Applicative order evaluation: starting with the innermost reducible expression (a redex).
- Normal order evaluation: starting with the outermost redex.
- If an expression has a normal form, normal order evaluation delivers it. Hence the name.
- For now you can imagine that Haskell uses normal order evaluation. A way to implement normal order evaluation is called *lazy evaluation*.

FUNCTIONS

MATHEMATICAL FUNCTIONS

- Mathematically, a function is a mapping between arguments and results.
 - A function $f :: A \rightarrow B$ maps each element in A to a unique element in B .
- In contrast, C “functions” are not mathematical functions:
 - `int y = 1; int f (x:int) { return ((y++) * x); }`
- Functions in Haskell have no such *side-effects*: (unconstrained) assignments, IO, etc.
- Why removing these useful features? We will talk about that later in this course.

CURRIED FUNCTIONS

- Consider again the function *smaller*:

smaller $:: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

smaller $x\ y = \text{if } x \leq y \text{ then } x \text{ else } y$

- We sometimes informally call it a function “taking two arguments”.
- Usage: *smaller* 3 4.
- Strictly speaking, however, *smaller* is a function returning a function. The type should be bracketed as $\text{Int} \rightarrow (\text{Int} \rightarrow \text{Int})$.

PRECEDENCE AND ASSOCIATION

- In a sense, all Haskell functions takes exactly one argument.
 - Such functions are often called *curried*.
- Type: $a \rightarrow b \rightarrow c = a \rightarrow (b \rightarrow c)$, not $(a \rightarrow b) \rightarrow c$.
- Application: $f\ x\ y = (f\ x)\ y$, not $f\ (x\ y)$.
 - *smaller 3 4* means *(smaller 3) 4*.
 - *square square 3* means *(square square) 3*, which results in a type error.
- Function application binds tighter than infix operators.
E.g. *square 3 + 4* means *(square 3) + 4*.

WHY CURRYING?

- It exposes more chances to reuse a function, since it can be partially applied.

twice $:: (a \rightarrow a) \rightarrow (a \rightarrow a)$

twice f x = *f (f x)*

quad $:: \text{Int} \rightarrow \text{Int}$

quad = *twice square*

- Try evaluating *quad 3*:

quad 3

= *twice square 3*

= *square (square 3)*

= ...

SECTIONING

- Infix operators are curried too. The operator $(+)$ may have type $Int \rightarrow Int \rightarrow Int$.
- Infix operator can be partially applied too.

$$(x \oplus) y = x \oplus y$$

$$(\oplus y) x = x \oplus y$$

- $(1 +) :: Int \rightarrow Int$ increments its argument by one.
- $(1.0 /) :: Float \rightarrow Float$ is the “reciprocal” function.
- $(/ 2.0) :: Float \rightarrow Float$ is the “halving” function.

INFIX AND PREFIX

- To use an infix operator in prefix position, surrounded it in parentheses. For example, `(+) 3 4` is equivalent to `3 + 4`.
- Surround an ordinary function by back-quotes (not quotes!) to put it in infix position. E.g. `3 'mod' 4` is the same as `mod 3 4`.

FUNCTION COMPOSITION

- Functions composition:

$$\begin{aligned}(\cdot) &:: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c) \\ (f \cdot g) \, x &= f \, (g \, x)\end{aligned}$$

- E.g. another way to write *quad*:

$$\begin{aligned}\textit{quad} &:: \textit{Int} \rightarrow \textit{Int} \\ \textit{quad} &= \textit{square} \cdot \textit{square}\end{aligned}$$

- Some important properties:
 - $\textit{id} \cdot f = f = f \cdot \textit{id}$, where $\textit{id} \, x = x$.
 - $(f \cdot g) \cdot h = f \cdot (g \cdot h)$.

GUARDED EQUATIONS

- Recall the definition:

$smaller :: Int \rightarrow Int \rightarrow Int$
 $smaller\ x\ y = \text{if } x \leq y \text{ then } x \text{ else } y$

- We can also write:

$smaller :: Int \rightarrow Int \rightarrow Int$
 $smaller\ x\ y \mid x \leq y = x$
 $\mid x > y = y$

- Helpful when there are many choices:

$signum :: Int \rightarrow Int$
 $signum\ x \mid x > 0 = 1$
 $\mid x = 0 = 0$
 $\mid x < 0 = -1$

λ EXPRESSIONS

- Since functions are first-class constructs, we can also construct functions in expressions.
- A λ expression denotes an anonymous function.
 - $\lambda x \rightarrow e$: a function with argument x and body e .
 - $\lambda x \rightarrow \lambda y \rightarrow e$ abbreviates to $\lambda x y \rightarrow e$.
 - In ASCII, we write λ as `\`
- Yet another way to define *smaller*:

smaller $:: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

smaller $= \lambda x y \rightarrow \text{if } x \leq y \text{ then } x \text{ else } y$

- Why λ s? Sometimes we may want to quickly define a function and use it only once.
- In fact, λ is a more primitive concept.

LOCAL DEFINITIONS

There are two ways to define local bindings in Haskell.

- **let**-expression:

```
f      :: Float → Float → Float
f x y = let a = (x + y)/2
          b = (x + y)/3
          in (a + 1) × (b + 2)
```

- **where**-clause:

```
f      :: Int → Int → Int
f x y | x ≤ 10 = x + a
      | x > 10 = x - a
      where a = square (y + 1)
```

- **let** can be used in expressions (e.g. $1 + (\mathbf{let..in..})$), while **where** qualifies multiple guarded equations.

TYPES

TYPES

- The universe of values is partitioned into collections, called *types*.
- Some basic types: *Int*, *Float*, *Bool*, *Char*...
- Type “constructors”: functions, lists, trees ...to be introduced later.
- Operations on values of a certain type might not make sense for other types. For example: *square square 3*.
- Strong typing: the type of a well-formed expression can be deduced from the constituents of the expression.
 - It helps you to detect errors.
 - More importantly, programmers may consider the types for the values being defined before considering the definition themselves, leading to clear and well-structured programs.

POLYMORPHIC TYPES

- Suppose $\text{square} :: \text{Int} \rightarrow \text{Int}$ and $\text{sqrt} :: \text{Int} \rightarrow \text{Float}$.
 - $\text{square} \cdot \text{square} :: \text{Int} \rightarrow \text{Int}$
 - $\text{sqrt} \cdot \text{square} :: \text{Int} \rightarrow \text{Float}$
- The $(\cdot)$ operator has different types in the two expressions:
 - $(\cdot) :: (\text{Int} \rightarrow \text{Int}) \rightarrow (\text{Int} \rightarrow \text{Int}) \rightarrow (\text{Int} \rightarrow \text{Int})$
 - $(\cdot) :: (\text{Int} \rightarrow \text{Float}) \rightarrow (\text{Int} \rightarrow \text{Int}) \rightarrow (\text{Int} \rightarrow \text{Float})$
- To allow $(\cdot)$ to be used in many situations, we introduce type variables and let its type be:
 $(b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c).$

SUMMARY SO FAR

- Functions are essential building blocks in a Haskell program. They can be applied, composed, passed as arguments, and returned as results.
- Types sometimes guide you through the design of a program.
- Equational reasoning: let the symbols do the work!

RECOMMENDED TEXTBOOKS

- *Introduction to Functional Programming using Haskell.*
My recommended book. Covers equational reasoning very well.
- *Programming in Haskell.* A thin but complete textbook.

ONLINE HASKELL TUTORIALS

- *Learn You a Haskell for Great Good!* , a nice tutorial with cute drawings!
- *Yet Another Haskell Tutorial*.
- *A Gentle Introduction to Haskell* by Paul Hudak, John Peterson, and Joseph H. Fasel: a bit old, but still worth a read.
- *Real World Haskell*. Freely available online. It assumes some basic knowledge of Haskell, however.