Programming Languages: Imperative Program Construction Practicals 6: Loop Constuction II

Shin-Cheng Mu

Autumn Term, 2024

- 1. Recall the maximum segment sum problem. What if we want to compute the maximum sum of *non-empty* segments?
 - (a) How would you write the specification? Does the specification still make sense with N being constrained only by $0 \leq N$?
 - (b) Derive a program solving the problem, assuming $-\infty$ to be the identity of (\uparrow).
- 2. Let n: Nat. Trying explaining why the following "splitting off" step is wrong by trying out the range calculation.

 $\begin{aligned} &\langle \Sigma i : n+1 \leqslant i < n+1 : f[i] \rangle \\ &= \begin{cases} \text{splitting off } i = n \\ &\langle \Sigma i : n+1 \leqslant i < n : f[i] \rangle + f[n] \end{cases}. \end{aligned}$

3. Derive a solution for:

con $N : Int\{N \ge 0\}$; $a : \operatorname{array} [0..N)$ of Intvar r : IntS $\{r = \langle \uparrow i, j : 0 \le i < j < N : a[i] - a[j] \rangle\}$.

4. Derive a solution for:

con $N : Int\{N \ge 1\}$; $a : \operatorname{array} [0..N)$ of Intvar r : IntS $\{r = \langle \#i, j : 0 \le i < j < N : a[i] \times a[j] \ge 0 \rangle\}$.

5. Consider again the maximum segment sum problem and its derivation in the handouts. Since the computation of *r* requires the value of the subterm $\langle \uparrow p : 0 \leq p \leq n+1 : sum p (n+1) \rangle$, some would think it makes more sense to use the following loop invariant $P_0 \land P_1 \land Q$, where

 $\begin{array}{l} P_0 \equiv r = \langle \uparrow p \; q : 0 \leqslant p \leqslant q \leqslant n : sum p \; q \rangle) \\ P_1 \equiv s = \langle \uparrow p : 0 \leqslant p \leqslant n + 1 : sum p \; (n + 1) \rangle \\ Q \equiv 0 \leqslant n \leqslant N . \end{array}$

What situation could you run into, if you try to construct a program using the invariant above? What if the array is non-empty, that is, $1 \leq N$?